

Mini-Camel

// identificatori
type ide = string ;;

// tipi
type tname = TInt | TBool | TClosure | TRecClosure | TUnBound ;;

// Espressioni astratte

type exp =
| EInt of int
| CstTrue
| CstFalse
| EString of string
| Den of ide

// Operatori binari int → int

| Sum of exp * exp
| Diff of exp * exp
| Prod of exp * exp
| Div of exp * exp

// Operatori int → bool

| IsZero of exp
| Eq of exp * exp
| LessThan of exp * exp
| GreaterThan of exp * exp

// Operatori su bool

| And of exp * exp
| Or of exp * exp
| Not of exp

// Controllo del flusso

| IfThenElse of exp * exp * exp

| Let of ide * exp * exp

| LetRec of ide * ide * exp * exp

| Fun of ide * exp

| Apply of exp * exp ;;

// Evaluation types

type evT =

| Int of int

| Bool of bool

| String of string

| Closure of ide * exp * evT env

| RecClosure of ide * ide * exp * evT env

| Unbound ;;

// Binding fra una stringa x e un valore primitivo V

let bind (s: evT env) (x: ide) (v: evT) =

function (i: ide) → if (i=x) then v else (s i);

// Get type

// Associa ad ogni valore il suo descrittore di tipo

let getType (x: evT): tname =

match x with

| Int(n) → TInt

| Bool(n) → TBool

| String(s) → TString

| Closure(i, e, en) → TClosure

| RecClosure(i, j, e, en) → TRecClosure

| Unbound → TUnbound

// Type checking

let typecheck $((x, y) : (\text{tname}^* \text{ evT})) =$

match x with

$| T \text{ Int} \rightarrow (\text{match } y \text{ with}$

$| \text{Int}(u) \rightarrow \text{true}$

$| - \rightarrow \text{false})$

$| T \text{ Bool} \rightarrow (\text{match } y \text{ with}$

$| \text{Bool}(u) \rightarrow \text{true}$

$| - \rightarrow \text{false})$

$| T \text{ String} \rightarrow (\text{match } y \text{ with}$

$| \text{String}(u) \rightarrow \text{true}$

$| - \rightarrow \text{false})$

$| T \text{ Closure} \rightarrow (\text{match } y \text{ with}$

$| \text{Closure}(i, e, en) \rightarrow \text{true}$

$| - \rightarrow \text{false})$

$| T \text{ RecClosure} \rightarrow (\text{match } y \text{ with}$

$| \text{RecClosure}(i, j, e, en) \rightarrow \text{true}$

$| - \rightarrow \text{false})$

$| T \text{ Unbound} \rightarrow (\text{match } y \text{ with}$

$| \text{Unbound} \rightarrow \text{true}$

$| - \rightarrow \text{false}) ;;$

// Interpretation

let rec eval $(e : \text{exp}) (s : \text{evT env}) : \text{evT} =$

match e with

$| E \text{ Int}(n) \rightarrow \text{Int}(n)$

$| \text{CstTrue} \rightarrow \text{Bool}(\text{true})$

$| \text{CstFalse} \rightarrow \text{Bool}(\text{false})$

$| E \text{ String} \rightarrow \text{String}(s)$

$| \text{Den}(i) \rightarrow (s i)$

- | Prod(e₁, e₂) → int-times ((eval e₁ s), (eval e₂ s))
- | Sum(e₁, e₂) → int-plus ((eval e₁ s), (eval e₂ s))
- | Diff(e₁, e₂) → int-sub ((eval e₁ s), (eval e₂ s))
- | Div(e₁, e₂) → int-div ((eval e₁ s), (eval e₂ s))
- | IsZero(e₁) → is-zero (eval e₁ s)
- | Eq(e₁, e₂) → int-eq ((eval e₁ s), (eval e₂ s))
- | LessThan(e₁, e₂) → less-than ((eval e₁ s), (eval e₂ s))
- | GreaterThan(e₁, e₂) → greater-than ((eval e₁ s), (eval e₂ s))
- | And(e₁, e₂) → bool-and ((eval e₁ s), (eval e₂ s))
- | Or(e₁, e₂) → bool-or ((eval e₁ s), (eval e₂ s))
- | Not(e₁) → bool-not (eval e₁ s)
- | IfThenElse(e₁, e₂, e₃) → let g = eval e₁ s in
 - (match (typecheck(TBool, g), g) with
 - | (true, Bool(true)) → eval e₂ s
 - | (true, Bool(false)) → eval e₃ s
 - | (-, -) → raise (RuntimeError "Wrong type"))
- | Let(i, e, ebody) → eval ebody (bind s i (eval e s))
- | Fun(arg, ebody) → Closure(arg, ebody, s)
- | LetRec(f, arg, fBody, letBody) →
 - let benv = bind s f (RecClosure(f, arg, fBody, s)) in
 - eval letBody benv
- | Apply(eF, eArg) → let fclosure = eval eF s in
 - (match fclosure with
 - | Closure(arg, fbody, fDecEnv) → let alld = eval eArg s in
 - let aenv = bind fDecEnv arg alld in
 - eval fbody aenv
- | RecClosure(f, arg, fbody, fDecEnv) → let alld = eval eArg s in
 - let xEnv = bind fDecEnv f closure in
 - let aenv = bind xEnv arg alld in
 - eval fbody aenv
- estendo l'ambiente statico*
- estendo xEnv con il passaggio dei parametri*
- | - → raise (RuntimeError "Wrong type"));

Semantica operazionale

ifthenelse (Regole di inferenza)

$$\frac{\Sigma \triangleright \text{cond} \Rightarrow \text{true} \quad \Sigma \triangleright e_1 \Rightarrow v_1}{\Sigma \triangleright \text{ifthenelse}(\text{cond}, e_1, e_2) \Rightarrow v_1}$$

$$\frac{\Sigma \triangleright \text{cond} \Rightarrow \text{false} \quad \Sigma \triangleright e_2 \Rightarrow v_2}{\Sigma \triangleright \text{ifthenelse}(\text{cond}, e_1, e_2) \Rightarrow v_2}$$

ifthenelse (Interprete)

ifthenelse (cond, e₁, e₂) →

let g = eval cond amb in
 match (typecheck ("bool", g), g) with
 | (true, Bool(true)) eval e₁ amb
 | (true, Bool(false)) eval e₂ amb
 | (-,-) → failwith "Non boolean guard"

Let (Regole di inferenza)

$$\frac{T \triangleright e_1 \Rightarrow v_1 \quad T[x=v_1] \triangleright e_2 \Rightarrow v_2}{T \triangleright \text{Let}(x, e_1, e_2) \Rightarrow v_2}$$

Let (Interprete)

let rec eval ((e: exp), (amb: envT env)) =
 match e with ...
 | Let (i, e, ebody) →
 eval ebody (bind amb i (eval e amb))

Function (Regole di inferenza)

Scoping statico

$$\frac{\begin{array}{l} \Gamma \triangleright \text{Fun}(x, e) \Rightarrow \text{Closure}(“x”, e, \Gamma) \\ \Gamma \triangleright \text{Var}(“f”) \Rightarrow \text{Closure}(“x”, \text{body}, \Gamma_{\text{closure}}) \\ \hline \Gamma \triangleright \text{arg} \Rightarrow \text{va} \quad \Gamma_{\text{closure}}[x=\text{va}] \triangleright \text{body} \Rightarrow v \\ \Gamma \triangleright \text{Apply}(\text{Den}(“f”), \text{arg}) \Rightarrow v \end{array}}{\Gamma \triangleright \text{Apply}(\text{Den}(“f”), \text{arg}) \Rightarrow v}$$

} Astrazione } Applicazione

$$\frac{\Gamma[f = \text{Closure}(“x”, e, \Gamma)] \triangleright e' \Rightarrow v'}{\Gamma \triangleright \text{Let}(“f”, \text{Fun}(x, e), e') \Rightarrow v'}$$

} Dichiarazione

Scoping dinamico

$$\frac{\begin{array}{l} \Gamma \triangleright \text{Fun}(“x”, e) \Rightarrow \text{Funval}(“x”, e) \\ \Gamma \triangleright \text{Den}(“f”) \Rightarrow \text{Funval}(“x”, e) \\ \hline \Gamma \triangleright \text{arg} \Rightarrow \text{va} \quad \Gamma[x=\text{va}] \triangleright e \Rightarrow v \\ \Gamma \triangleright \text{Apply}(\text{Den}(“f”), \text{arg}) \Rightarrow v \end{array}}{\Gamma \triangleright \text{Apply}(\text{Den}(“f”), \text{arg}) \Rightarrow v}$$

} Astrazione } Applicazione

Function (Interprete) Scoping statico

```
let rec eval ((e:exp)(amb:env env)) =  
  match e with ...  
  | Fun(i,a) → Closure(i,a,amb)  
  | Apply(Den(f), eArg) →  
    let fclosure = lookup amb f in  
    (match fclosure with  
     | Closure(arg, fbody, fDecEnv) →  
       let aVal = eval eArg amb in  
       let aEnv = bind fDecEnv arg aVal in  
       eval fbody aEnv  
     | _ → failwith "Non functional value"  
  | Apply(-,-) → failwith "Application: not first order function";;
```

Scoping dinamico

type env = Int of int

| Bool of bool

| Unbound

| Funval of efun

and efun = ide * exp

let rec eval ((e: exp) (amb: envT env)) =
match e with ...
| Fun (arg, ebody) → Funval (arg, ebody)
| Apply (Den(f), eArg) →
 let fval = lookup amb f in
 (match fval with
 | Funval (arg, fbody) →
 let aVal = eval eArg amb in
 let aenv = bind amb arg aVal in
 eval fbody aenv
 | _ → failwith "Non functional value"
 | Apply (-, -) → failwith "Application: not first order function";;