

Esami

1) $T \triangleright \text{BothFun}(x, e_1, e_2) \Rightarrow \text{BothClosure}(x, e_1, e_2, T)$

$T \triangleright e_1 \Rightarrow \text{BothClosure}(\text{arg}, \text{body}_1, \text{body}_2, T'_{\text{fdec}})$ $T \triangleright e_2 \Rightarrow \text{aval}$

$T'_{\text{fdec}}[\text{arg} = \text{aval}] \triangleright \text{body}_1 \Rightarrow v_1$ $T'_{\text{fdec}}[\text{arg} = \text{aval}] \triangleright \text{body}_2 \Rightarrow v_2$
 $T \triangleright \text{Apply}(e_1, e_2) \Rightarrow (v_1, v_2)$

type exp = ...

| BothFun of ide * exp * exp

type evT = ...

| BothClosure of ide * exp * exp * evT env

| Pair of exp * exp

let rec eval e s =

match e with ...

| BothFun(arg, body1, body2) →

 BothClosure(arg, body1, body2, s)

| Apply(e1, e2) →

 let fclosure = eval e1 s in

 (match fclosure with ...

 | BothClosure(arg, body1, body2, fDecEnv) →

 let aval = eval e2 s in

 let aenv = bind fDecEnv arg aval in

 Pair(eval body1 aenv, eval body2 aenv))

2) type tname = ... | T Coda

type exp = ...

| CodaLimitata of exp

| Insert of exp * exp

| Remove of exp * exp

| Peek of exp

type evT = ...

| Coda of evT list * int

let typecheck (x, y) =

match (x, y) with

| (T Coda, Coda(lis, m)) → true

| (-, -) → false

let rec eval e s =

match e with ...

| CodaLimitata e1 → let size = eval e1 s in

if typecheck(TInt, size) then

(match size with

| TInt(m) → Coda([], m)

| _ → failwith "Error")

| Insert (elem, e1) → let coda = eval e1 s in

(match coda with

| Coda(lis, size) →

if List.length lis < size then

let x = eval elem s in

Coda(lis @ [x], size)

else failwith "Coda piena"

$| _ \rightarrow \text{failwith "Error"}$)
 $| \text{Remove}(\text{elem}, e_1) \rightarrow \text{let coda} = \text{eval } e_1 \text{ s in}$
 (match coda with)
 $| \text{Coda}([], \text{size}) \rightarrow$
 $\quad \text{failwith "Coda vuota"}$
 $| \text{Coda}(x::\text{lis}, \text{size}) \rightarrow$
 $\quad \text{Coda}(\text{lis}, \text{size})$
 $| _ \rightarrow \text{failwith "Error"}$)

$| \text{Peek } e_1 \rightarrow \text{let coda} = \text{eval } e_1 \text{ s in}$
 (match coda with)
 $| \text{Coda}([], \text{size}) \rightarrow$

Alternativa

$// \text{Coda}(x::\text{lis}, \text{size}) \rightarrow x$
 $| \text{Coda}(\text{lis}, \text{size}) \rightarrow$
 $\quad \text{List. hd lis}$
 $| _ \rightarrow \text{failwith "Error"}$)

3) Pixel $p := \langle r, g, b \rangle$ dove $r, g, b \in \{0, \dots, 255\}$

Identificatori $I := \dots$

Esempi $e := I | p | \text{lighten } e | \text{darken } e | \text{let } I = e_1 \text{ in } e_2$

type pixel = int * int * int

type evT = Unbound | Epix of pixel

type ide = string

type exp = Ide of ide | Pix of pixel
 $| \text{Light of exp} | \text{Dark of exp}$
 $| \text{Let of ide * exp * exp}$

let check $a = a > -1 \ \&\& a < 256$

let inc $a = \text{match } a \text{ with}$

$$| 255 \rightarrow 255$$

$$| n \rightarrow n+1$$

let dec $a = \text{match } a \text{ with}$

$$| 0 \rightarrow 0$$

$$| n \rightarrow n-1$$

let increase $v = \text{match } v \text{ with}$

$$| Epix(a, b, c) \rightarrow Epix(\text{inc } a, \text{inc } b, \text{inc } c)$$

let decrease $v = \text{match } v \text{ with}$

$$| Epix(a, b, c) \rightarrow Epix(\text{dec } a, \text{dec } b, \text{dec } c)$$

let rec eval $e\ s =$

$\text{match } e \text{ with} \dots$

$$| \text{Id}(i) \rightarrow s\ i$$

$$| \text{Pix}(a, b, c) \rightarrow \text{if check } a \ \&\& \text{check } b \ \&\& \text{check } c \text{ then} \\ \quad Epix(a, b, c) \text{ else UnBound}$$

$$| \text{Light } e_1 \rightarrow \text{let } v = \text{eval } e_1\ s \text{ in increase } v$$

$$| \text{Dark } e_1 \rightarrow \text{let } v = \text{eval } e_1\ s \text{ in decrease } v$$

$$| \text{Let}(i, e_1, e_2) \rightarrow \text{let } v = \text{eval } e_1\ s \text{ in let } s_1 = \text{bind}$$

4) for-each (lista-intei; funzione)

type exp = ...

| ForEach of exp list * exp

let rec eval e s =

match e with ...

| ForEach e1 e2 → (let lis = evalSeq e1 s in

let f = eval e2 s in

match (lis, f) with

| (lst, Closure (i, e, s1)) →

Int (evalList lst (i, e, s1) s)

| (-, -) → failwith "Error"

and evalSeq lst s =

match lst with

| [] → []

| e :: lst' → let x = eval e s in

(match (typecheck (TInt, x), x) with

| (true, Int v) → v :: (evalSeq lst' s)

| (false, _) → failwith "Error"

and evalList lst funB s =

match (lst, funB) with

| ([], _) → 0

| (v :: lst', (args, fBody, fDecEnv)) →

match (eval fBody (bind fDecEnv args (Int v))) with

| Int v → v + (evalList lst' funB s);

5)

$$\frac{T \triangleright e \Rightarrow v \quad T' = T[f \mapsto \text{MyClosure}(x, v, e_1, T)] \quad T' \triangleright e_2 \Rightarrow v_1]}{T \triangleright \text{function } f(x=e) = e_1 \text{ in } e_2 \Rightarrow v_1}$$

$T \triangleright \text{function } f(x=e) = e_1 \text{ in } e_2 \Rightarrow v_1$

$$\frac{T(f) = \text{MyClosure}(x, v, e, T_{f\text{Decl}}) \quad T \triangleright e\text{arg} \Rightarrow v\text{arg} \quad T_{f\text{Decl}}[x \mapsto v\text{arg}] \triangleright e \Rightarrow v_1}{T \triangleright f(e\text{arg}) \Rightarrow v_1}$$

$$\frac{T(f) = \text{MyClosure}(x, v, e, T_{f\text{Decl}}) \quad T_{f\text{Decl}}[x \mapsto v] \triangleright e \Rightarrow v_1}{T \triangleright f() \Rightarrow v_1}$$

type exp = ...

- | MyFun of ide * ide * exp * exp * exp
- | MyApply of ide * (exp option)

type evT = ...

- | MyClosure of ide * evT * exp * evT env

let rec eval e s =

match e with ...

| MyFun(f, x, edef, ebody, e) →

let vdef = eval edef s in

let closure = MyClosure(x, vdef, ebody, s) in

let s1 = bind s f closure in

eval e s1

| MyApply(f, arg) → let closure = s f in

(match closure with

| MyClosure(x, vdef, ebody, s1) →

(match arg with

| Some e →

let $s_2 = \text{bind } s_1 \times (\text{eval arg } s_1) \text{ in}$
 $\text{eval ebody } s_2$

| None \rightarrow

 let $s_2 = \text{bind } s_1 \times \text{vdef } i \text{ in}$
 $\text{eval ebody } s_2$)

| _ \rightarrow failwith "Error");;

⑥)

$\frac{\Gamma \triangleright \text{cond1} \Rightarrow \text{true} \quad \Gamma \triangleright \text{cond2} \Rightarrow \text{true} \quad \Gamma \triangleright e_1 \Rightarrow v_1}{\Gamma \triangleright \text{on } (\text{cond1}, e_1, \text{cond2}, e_2) \Rightarrow v_1}$

$\frac{\Gamma \triangleright \text{cond1} \Rightarrow \text{true} \quad \Gamma \triangleright \text{cond2} \Rightarrow \text{false} \quad \Gamma \triangleright e_1 \Rightarrow v_1}{\Gamma \triangleright \text{on } (\text{cond1}, e_1, \text{cond2}, e_2) \Rightarrow v_1}$

$\frac{\Gamma \triangleright \text{cond1} \Rightarrow \text{false} \quad \Gamma \triangleright \text{cond2} \Rightarrow \text{true} \quad \Gamma \triangleright e_2 \Rightarrow v_2}{\Gamma \triangleright \text{on } (\text{cond1}, e_1, \text{cond2}, e_2) \Rightarrow v_2}$

type exp = ...

| on of exp * exp * exp * exp

let rec eval e s =

match e with ...

| on (cond1, e1, cond2, e2) \rightarrow

 let g1 = eval cond1 s in

 (match (typecheck (TBool, g1), g1) with

 | (true, Bool(true)) \rightarrow let g2 = eval cond2 s in

 (match (typecheck (TBool, g2), g2) with

 | (true, Bool(true)) \rightarrow eval e1 s

 | (true, Bool(false)) \rightarrow eval e1 s

 | (_, _) \rightarrow failwith "Type error")

| (true, $\text{Bool}(\text{false})$) \rightarrow let $g_2 = \text{eval cond.2 } s$ in
(match (typecheck(T_{Bool} , g_2), g_2) with
| (true, $\text{Bool}(\text{true})$) \rightarrow eval e_2 s
| (true, $\text{Bool}(\text{false})$) \rightarrow raise "False guards"
| (-, -) \rightarrow failwith "type error")

7)

type exp = ...

| use of $\text{exp}^* \text{exp}^* \text{exp}^* \text{exp}$

let rec eval e $s =$

match e with ...

| use ($e_1, e_2, e_3, \text{cond}$) \rightarrow

let $g = \text{eval cond } s$ in

(match (typecheck(T_{Bool} , g), g) with

| (true, $\text{Bool}(\text{true})$) \rightarrow

let $s_1 = \text{bind } s \times (\text{eval } e_1 \ s)$ in

eval e_3 s_1

| (true, $\text{Bool}(\text{false})$) \rightarrow

let $s_1 = \text{bind } s \times (\text{eval } e_2 \ s)$ in

eval e_3 s_1

| (-, -) \rightarrow failwith "type error")

8)

$$\Gamma \triangleright F u (x, e) \rightarrow F u \text{ Closure} ("x", e, \Gamma)$$

$$\Gamma \triangleright V a r ("F u") \rightarrow F u \text{ Closure} ("x", \text{body}, \Gamma_{\text{free}})$$

$$\frac{\Gamma \triangleright \text{arg} \Rightarrow v_a \quad \Gamma_{\text{free}} [x=v_a] \triangleright \text{body} \Rightarrow v}{\Gamma \triangleright F u \text{ Apply} (\text{Den} ("F u"), \text{arg}) \Rightarrow v}$$

type exp = ...

| F u of ide * exp * exp list

| F u Apply of ide * exp

type evT = ...

| F u Closure of ide * exp * exp list * evT env

let rec eval e s =

match e with

| F u (arg, fbody, lis) \rightarrow F u Closure (arg, fbody, lis, f DecEnv)| Apply (eF, eArg) \rightarrow

let fclosure = eval eF s in

(match fclosure with ...)

| F u Closure (arg, fbody, lis, f DecEnv) \rightarrow

let aVal = eval eArg s in

let rec search x lst =

let typeName = getType x in

match lst with

| [] \rightarrow failwith "Elemento non presente"| y::l \rightarrow (match typecheck (typeName, y) with| true \rightarrow if x=y then y else search x l| false \rightarrow search x l)

in

let aenv = bind fDecEnv arg (List.map (fun x → eval x s)
lis) in

eval fbody aenv

|_ → failwith "Error"

9) Posizione $p := 0, \dots, 100$

Identifieri $I := \dots$

Espressioni $e := I | p | \text{moveLeft } e | \text{moveRight } e | \text{let } I = e_1 \text{ in } e_2$

type ide = string

exception RuntimeError of string

type evT = ...

| Unbound

| Pos of int

| String of string

type exp = ...

| I of ide

| Epos of int

| MoveLeft of exp

| MoveRight of exp

| Let of ide * exp * exp

type 't env = ide → 't // Ambiente polimorfo

let emptyEnv = function x → Unbound // Ambiente vuoto

```

let bind (s: evT env) (i: idt) (v: evT) =
  function (x: idt) → if i = x then v else (s i)

let typecheck ((x, y): (tname * evT)) =
  match x with
  | TInt → (match y with
    | Pos(m) → true
    | _ → raise (RuntimeException "Wrong type"))
  | TString → (match y with
    | String(s) → true
    | _ → raise (RuntimeException "Wrong type"))

```

```

let moveL x =
  match (typecheck (TInt, x), x) with
  | (true, Pos(m)) → if m = 0 then Pos(0) else Pos(m - 1)
  | _ → failwith "Error"

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let moveR x =
  match (typecheck (TInt, x), x) with
  | (true, Pos(m)) → if m = 100 then Pos(100) else Pos(m + 1)
  | _ → failwith "Error"

```

```

let rec eval e s =
  match e with
  | I(u) → (s u)
  | EPos(u) → Pos(u)
  | MoveLeft(e1) → let p = eval e1 s in moveL p
  | MoveRight(e1) → let p = eval e1 s in moveR p
  | Let(i, e1, body) → eval body (bind s i (eval e1 s))

```

